

Strongly-MDS Convolutional Codes

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Abstract—Maximum-distance separable (MDS) convolutional codes have the property that their free distance is maximal among all codes of the same rate and the same degree. In this paper, a class of MDS convolutional codes is introduced whose column distances reach the generalized Singleton bound at the earliest possible instant. Such codes are called strongly-MDS convolutional codes. They also have a maximum or near-maximum distance profile. The extended row distances of these codes will also be discussed briefly.

Index Terms—Column distances, convolutional codes, extended row distances, maximum-distance separable (MDS) codes, super-regular matrices, unit memory codes.

I. INTRODUCTION

IN comparison to the literature on linear block codes, there exist only relatively few algebraic constructions of convolutional codes having some good designed distance. There are even fewer algebraic decoding algorithms which are capable of exploiting the algebraic structure of the code.

A large part of the literature on convolutional codes studies codes over the binary field. Codes are then typically presented by trellis and state diagrams and the decoding algorithm of choice is the Viterbi algorithm. The reader is referred to the standard books [1], [2] or the more recent articles [3]–[5] where also further references to the literature can be found.

In the early 1970s, there were some important constructions done for convolutional codes over large alphabets and we would like to mention the papers [6]–[9] and the monograph by Piret [10]. In [7], [8], Justesen and Hughes study the question on how large the free distance of a convolutional code over a large alphabet can be. In [11], the authors of the present paper derived a generalized Singleton bound and they define a convolutional code to be maximum-distance separable (MDS) if its free distance reaches this upper bound. Using a construction idea due to Justesen [6], they provided in [12] a first concrete construction of MDS convolutional codes for all rates and degrees.

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In [13], the class of cyclic convolutional codes, as first introduced by Piret [14] and Roos [15], has been studied. It turned out that many of the constructed codes were MDS and/or were codes over large alphabets with excellent distance profile. More recently, Goppa convolutional codes have been introduced in [16] and many of these codes have excellent distances as well.

In this paper, we will consider the problem of constructing codes with rapid growth of their column distances. These distance parameters have been introduced by Costello in [17] and have been studied by many authors since. See, e.g., [1], [18], and the references therein.

Algorithmic searches for codes with large column distances have been carried out for instance in [19], [20]. We will now attack the question of convolutional codes with large column distances from the theoretical point of view. First, we will discuss how big these distances can possibly be, thereby introducing a new class of convolutional codes which we call strongly-MDS convolutional codes. These are codes with a column distance profile such that the free distance, i.e., the generalized Singleton bound, is reached at the earliest possible stage. The main part of the paper will be about existence and construction of such codes. At the end of the paper, we will also use our methods in order to briefly discuss the extended row distances of these codes in case they have unit memory. As opposed to the column distances, the extended row distances also grow beyond the free distance and thus contain some additional information about the performance of the code.

Convolutional codes over large alphabets are interesting both from a purely theoretical as well as from an applications point of view. On the theory side, the following questions arise naturally: How large can be the free distance of a convolutional code of some fixed rate and fixed degree? How to construct codes which achieve a maximal free distance? How good can be the column and row distance profile of these codes?

From a practical point of view, we can identify a convolutional code over a finite alphabet with a finite linear state machine (LFSM) having redundancy and which is capable of correcting processing errors. In a series of recent papers [21]–[23], Hadjicostis and Verghese showed how to error protect a given LFSM with a larger redundant LFSM capable of detecting and correcting state transition errors. The detection and correction of errors is done using nonconcurrent measurements of the state vectors of the redundant LFSM. The construction of the redundant system boils down to the construction of convolutional codes with good free distance over an alphabet which is in general nonbinary.

Let \mathbb{F} be any finite field and let $\mathbb{F}[D]$ be the polynomial ring. For the definition of a convolutional code we take a module theoretic point of view [24], [25].

Definition 1.1: A convolutional code of rate k/n is a submodule $\mathcal{C} \subseteq \mathbb{F}[D]^n$ of rank k such that there exists a $k \times n$

polynomial encoder matrix $G \in \mathbb{F}[D]^{k \times n}$ which is *basic*, i.e., G has a polynomial right inverse, and which is *minimal*, i.e., the sum of the row degrees of G attains the minimal possible value, such that

$$\mathcal{C} := \{uG \mid u \in \mathbb{F}[D]^k\} \subseteq \mathbb{F}[D]^n.$$

We define the *degree* of \mathcal{C} as the sum of the row degrees of one, and hence any, minimal basic encoder.

In the sequel, we will adopt the notation used by McEliece [26, p. 1082] and call a convolutional code of rate k/n and degree δ an (n, k, δ) code. Every (n, k, δ) code \mathcal{C} can be presented in terms of a parity-check matrix $H \in \mathbb{F}[D]^{(n-k) \times n}$, i. e.,

$$\mathcal{C} = \{v \in \mathbb{F}[D]^n \mid vH^T = 0\}.$$

It is clear that we can choose H to be basic and minimal and we will do so throughout the paper. Notice also that $GH^T = 0$ for any generator matrix G of \mathcal{C} .

The *weight* of a vector $v = \sum_{j=0}^L v_j D^j \in \mathbb{F}[D]^n$ is defined as

$$\text{wt}(v) := \sum_{j=0}^L \text{wt}(v_j) \in \mathbb{N}_0$$

where $\text{wt}(v_j)$ denotes the Hamming weight of $v_j \in \mathbb{F}^n$. The *free distance* of the code $\mathcal{C} \subset \mathbb{F}[D]^n$ is given as

$$d_{\text{free}} := \min\{\text{wt}(v) \mid v \in \mathcal{C}, v \neq 0\}.$$

Since we assume that $G \in \mathbb{F}[D]^{k \times n}$ is minimal and basic, we also have

$$d_{\text{free}} = \min\{\text{wt}(v) \mid v = uG \text{ for some } u \in \mathbb{F}[D]^k \setminus \{0\}, u_0 \neq 0\}.$$

In Theorem 2.6, we will recall from the paper [11] an upper bound on d_{free} based on the parameters (n, k, δ) . It generalizes the Singleton bound from block code theory and will play a central role in our paper.

The paper is structured as follows: In Section II, we review notions from convolutional coding theory such as column distances, the generalized Singleton bound and we introduce the important concepts for this paper, namely, the property of being strongly-MDS and having a maximum distance profile. In Section III, we show the existence of strongly-MDS codes of rate k/n with $n-k \mid \delta$. In order to do so, we introduce the interesting concept of a *superregular matrix* which could be of independent interest. In Section IV, we illustrate these concepts through a series of examples. In Section V, we investigate to what extent properties of MDS, strongly-MDS, and maximum distance profile carry over to the dual code. The main result of this section states that a code has a maximum distance profile if and only if its dual has this property. This allows us to show that for certain specific parameters, a code is strongly-MDS if and only if its dual is strongly-MDS. Finally, in Section VI, we will derive a lower bound for the extended row distances.

II. STRONGLY-MDS CODES AND CODES WITH MAXIMUM DISTANCE PROFILE

Let $\mathcal{C} \subseteq \mathbb{F}[D]^n$ be an (n, k, δ) code with a basic minimal generator matrix

$$G = \sum_{j=0}^{\nu} G_j D^j \in \mathbb{F}[D]^{k \times n}, G_j \in \mathbb{F}^{k \times n}, G_{\nu} \neq 0 \quad (2.1)$$

and a basic parity-check matrix

$$H = \sum_{j=0}^{\mu} H_j D^j \in \mathbb{F}[D]^{(n-k) \times n}, H_j \in \mathbb{F}^{(n-k) \times n}, H_{\mu} \neq 0. \quad (2.2)$$

Notice that ν is the memory of the code. For every $j \in \mathbb{N}_0$, the truncated sliding generator matrices $G_j^c \in \mathbb{F}^{(j+1)k \times (j+1)n}$ and parity-check matrices $H_j^c \in \mathbb{F}^{(j+1)(n-k) \times (j+1)n}$ are given by

$$G_j^c := \begin{bmatrix} G_0 & G_1 & \dots & G_j \\ & G_0 & \dots & G_{j-1} \\ & & \ddots & \vdots \\ & & & G_0 \end{bmatrix} \quad (2.3)$$

$$H_j^c := \begin{bmatrix} H_0 & & & \\ H_1 & H_0 & & \\ \vdots & \vdots & \ddots & \\ H_j & H_{j-1} & \dots & H_0 \end{bmatrix}$$

where we let $G_j = 0$ (resp., $H_j = 0$) whenever $j > \nu$ (resp., $j > \mu$); see also [1, p. 110]. The identity $GH^T = 0$ and the full ranks of G_0 and H_0 immediately imply the full rank of the sliding matrices as well as the identities

$$\{uG_j^c \mid u \in \mathbb{F}^{(j+1)k}\} = \{v \in \mathbb{F}^{(j+1)n} \mid v(H_j^c)^T = 0\}$$

for all $j \in \mathbb{N}_0$. Using these equations and following [1, pp. 110], the j th *column distance* of \mathcal{C} is given as (2.4)–(2.5) at the bottom of the page. The column distances are invariants of the code, i.e., they do not depend on the choice of the generator matrix (see [1, Sec. 3.1]), and satisfy

$$d_0^c \leq d_1^c \leq d_2^c \dots \leq \lim_{j \rightarrow \infty} d_j^c = d_{\text{free}}. \quad (2.6)$$

The $(\nu + 1)$ -tuple of numbers $(d_0^c, d_1^c, d_2^c, \dots, d_{\nu}^c)$, where ν is the memory, is called the *column distance profile* of the code [1, p. 112].

Equation (2.5) immediately implies the following simple fact used several times in the paper.

Proposition 2.1: Let $d \in \mathbb{N}$. Then the following properties are equivalent:

- $d_j^c = d$;
- none of the first n columns of H_j^c is contained in the span of any other $d - 2$ columns and one of the first n columns of H_j^c is in the span of some other $d - 1$ columns of that matrix.

$$d_j^c = \min\{\text{wt}((u_0, \dots, u_j)G_j^c) \mid u_i \in \mathbb{F}^k, u_0 \neq 0\} \quad (2.4)$$

$$= \min\{\text{wt}(\hat{v}) \mid \hat{v} = (\hat{v}_0, \dots, \hat{v}_j) \in \mathbb{F}^{(j+1)n}, \hat{v}(H_j^c)^T = 0, \hat{v}_0 \neq 0\}. \quad (2.5)$$

We leave it to the reader to verify the equivalence of the statements. The following upper bound on the column distances is an immediate consequence of the previous result, along with the fact that H_j^c has full row rank.

Proposition 2.2: For every $j \in \mathbb{N}_0$ we have

$$d_j^c \leq (n-k)(j+1) + 1.$$

This observation has already been made in the context of systematic convolutional codes in [8] and in [27]. The next proposition shows that maximality of d_j^c implies maximality of the preceding column distances.

Corollary 2.3: If $d_j^c = (n-k)(j+1) + 1$ for some $j \in \mathbb{N}_0$, then $d_i^c = (n-k)(i+1) + 1$ for all $i \leq j$.

Proof: It suffices to prove the assertion for $i = j-1$. In order to do so notice that

$$H_j^c = \left[\begin{array}{c|c} H_{j-1}^c & \begin{matrix} 0 \\ \vdots \\ 0 \end{matrix} \\ \hline H_j & H_0 \end{array} \right]$$

and assume that one of the first n columns of H_{j-1}^c is in the span of some other $(n-k)j-1$ columns. Then $\text{rank } H_0 = n-k$ implies that one of the first n columns of H_j^c is in the span of some other $(n-k)j-1+n-k = (n-k)(j+1)-1$ columns of H_j^c . But this is a contradiction to the optimality of d_j^c by Proposition 2.1. \square

The following characterizations of the j th column distance being maximal will be useful later on for constructing strongly-MDS codes and also when considering duality in Section V.

Theorem 2.4: Let G_j^c and H_j^c be as in (2.3). Then the following are equivalent:

- i) $d_j^c = (n-k)(j+1) + 1$;
- ii) every $(j+1)k \times (j+1)k$ full-size minor of G_j^c formed from the columns with indices $1 \leq t_1 < \dots < t_{(j+1)k}$, where $t_{sk+1} > sn$ for $s = 1, \dots, j$, is nonzero;
- iii) every $(j+1)(n-k) \times (j+1)(n-k)$ full-size minor of H_j^c formed from the columns with indices $1 \leq r_1 < \dots < r_{(j+1)(n-k)}$, where $r_{s(n-k)} \leq sn$ for $s = 1, \dots, j$, is nonzero.

Notice that the index condition in part ii) simply says that for each s , the minors under consideration have at most sk columns out of the first sn columns of G_j^c . Clearly, all other full-size minors of G_j^c are singular. The proof of this purely matrix theoretical result will be given in Appendix A. The next corollary will provide a link of our results to the existing literature later in this section. It will also be helpful for the construction of codes with maximum column distances in the next section.

Corollary 2.5: Let $G = [I, P]$ be a systematic matrix of memory ν and let $P = \sum_{i=0}^{\nu} P_i D^i \in \mathbb{F}[D]^{k \times (n-k)}$ where $P_\nu \neq 0$. Then the following are equivalent for $j = 0, \dots, \nu$:

- i) $d_j^c = (n-k)(j+1) + 1$;

- ii) each $i \times i$ -submatrix of

$$\hat{P} := \begin{bmatrix} P_0 & P_1 & \dots & P_j \\ & P_0 & \dots & P_{j-1} \\ & & \ddots & \vdots \\ & & & P_0 \end{bmatrix} \in \mathbb{F}^{(j+1)k \times (j+1)(n-k)}$$

that does not contain an $s \times t$ -zero block for some s, t such that $s+t \geq i+1$ is nonsingular. If one of these conditions is satisfied, then all entries of the matrices P_0, \dots, P_j are nonzero.

This result appeared already in [27]. However, since the paper is not easily accessible and no detailed proof is provided, we added a proof in Appendix A. We will come back to the relation of our work with [27] at the end of this section.

In the sequel, we will relate the upper bound for the column distances to an upper bound for the free distance of the code. The maximum possible value for the free distance of a convolutional code over any field has been established in [11]. Therein the following has been shown.

Theorem 2.6: The free distance of an (n, k, δ) code satisfies

$$d_{\text{free}} \leq (n-k) \left(\left\lfloor \frac{\delta}{k} \right\rfloor + 1 \right) + \delta + 1. \quad (2.7)$$

Notice that if $\delta = 0$, the number on the right-hand side of (2.7) reduces to the usual Singleton bound $n-k+1$ for an (n, k) block code. Therefore, we call this number the *generalized Singleton bound* and codes whose distance attains this bound will be called *MDS codes*. It has been shown in [11] that for every set of parameters (n, k, δ) and every prime number p there exists a suitably large finite field \mathbb{F} of characteristic p and an MDS code with parameters (n, k, δ) over \mathbb{F} . The proof is based on techniques from algebraic geometry and is nonconstructive. In [12], a construction of MDS codes with parameters (n, k, δ) was given for suitably large fields of characteristic coprime with n .

In the sequel, we will strengthen the MDS property by requiring that the generalized Singleton bound is attained by the earliest column distance possible. This will lead us to the notion of a strongly-MDS code.

Proposition 2.7: Suppose \mathcal{C} be an (n, k, δ) MDS code with column distances d_j^c , $j \in \mathbb{N}_0$ and free distance d_{free} . Let $M := \min\{j \in \mathbb{N}_0 \mid d_j^c = d_{\text{free}}\}$. Then

$$M \geq \left\lfloor \frac{\delta}{k} \right\rfloor + \left\lceil \frac{\delta}{n-k} \right\rceil.$$

Proof: From Proposition 2.2 we get

$$d_{\text{free}} = (n-k) \left(\left\lfloor \frac{\delta}{k} \right\rfloor + 1 \right) + \delta + 1 = d_M^c \leq (n-k)(M+1) + 1. \quad (2.8)$$

This yields the assertion. \square

The proof also shows that in the case $j > \left\lfloor \frac{\delta}{k} \right\rfloor + \left\lceil \frac{\delta}{n-k} \right\rceil$, the column distance d_j^c never attains the upper bound $(n-k)(j+1) + 1$ of Proposition 2.2, see also (2.6).

Definition 2.8: An (n, k, δ) code with column distances d_j^c , $j \in \mathbb{N}_0$, is called *strongly-MDS*, if

$$d_M^c = (n-k) \left(\left\lfloor \frac{\delta}{k} \right\rfloor + 1 \right) + \delta + 1 \text{ for } M = \left\lfloor \frac{\delta}{k} \right\rfloor + \left\lceil \frac{\delta}{n-k} \right\rceil.$$

Because of (2.6) strongly-MDS codes are in particular MDS codes.

In the case where $(n - k) \mid \delta$, the strong-MDS property implies that d_M^c attains the upper bound $(n - k)(M + 1) + 1$, see Proposition 2.1. Hence, Corollary 2.3 shows that in this case *all* column distances attain their optimum value. In other words, the column distances are

$$(d_0^c, d_1^c, d_2^c, \dots) \\ = (n - k + 1, 2(n - k) + 1, \dots, M(n - k) + 1, S, S, \dots) \quad (2.9)$$

where $S = (M + 1)(n - k) + 1$. If $(n - k) \nmid \delta$, we always have $d_M^c < (n - k)(M + 1) + 1$ as can be seen from (2.8). In this case, nothing can be concluded for the previous column distances. In order to also incorporate optimality of d_0^c, \dots, d_{M-1}^c in this general case we pose the following definition.

Definition 2.9: Let

$$L := \left\lfloor \frac{\delta}{k} \right\rfloor + \left\lfloor \frac{\delta}{n - k} \right\rfloor. \quad (2.10)$$

An (n, k, δ) code with column distances d_j^c , $j \in \mathbb{N}_0$ is said to have a maximum distance profile if

$$d_j^c = (n - k)(j + 1) + 1, \quad \text{for } j = 1, \dots, L.$$

Using the notation of Definition 2.8 we have

$$L = \begin{cases} M, & \text{if } (n - k) \mid \delta \\ M - 1, & \text{otherwise.} \end{cases} \quad (2.11)$$

Obviously, a strongly-MDS code with maximum-distance profile satisfies (2.9) for $S = (n - k)(\lfloor \frac{\delta}{k} \rfloor + 1) + \delta + 1$. Furthermore, we obtain that if $n - k$ divides δ then an (n, k, δ) code has maximum distance profile if and only if it is strongly-MDS. If $n - k$ does not divide δ , then neither property implies the other one as is verified by Examples in 2.12 later in the paper.

An immediate consequence of Corollary 2.3 is as follows.

Remark 2.10: An (n, k, δ) code has a maximum distance profile if and only if the L th column distance satisfies

$$d_L^c = (n - k)(L + 1) + 1.$$

Remark 2.11: The concept is clearly related to the notion of optimum distance profile (ODP), see [1, p. 112]. For ODP it is required that the column distances are maximal up to the memory ν . Hence, if $\nu \leq L$ then a code with maximum distance profile is always ODP. In general, one expects a good code to have generic Forney indices, i.e., the indices attain only the two values $\lceil \frac{\delta}{k} \rceil$ and $\lfloor \frac{\delta}{k} \rfloor$. McEliece [26, Corollary 4.3] calls such codes *compact codes*. It has been shown in [11] that an MDS code always has generic indices. Of course, if the indices are generic then $\nu = \lceil \frac{\delta}{k} \rceil$ and thus, $\nu \leq L + 1$.

The notion of ODP seems also to be dependent on the base field which is usually assumed to be the binary field. A code with maximum distance profile, in general, does not exist over the binary field and it can only exist for sufficiently large base fields. This is similar to the situation of MDS block codes. Such codes are known to exist as soon as the field size of \mathbb{F} is larger than the block length n .

We close this section by relating our work to previous results in the literature. As indicated earlier, the papers [8], [27], [28]

deal with a notion closely related to the bound given in Proposition 2.2. In all these papers (n, k) convolutional codes with a systematic generator matrix of memory m are considered and such codes are called MDS if $d_m^c = (n - k)(m + 1) + 1$. In order to avoid confusion, we will for the rest of this section call codes that satisfy $d_m^c = (n - k)(m + 1) + 1$, where m is the memory, m -MDS codes. It is easy to see that the number $(n - k)(m + 1) + 1$ is the maximum possible value for the free distance of a *systematic* convolutional code with these parameters. In the papers [27], [28] this property has been characterized by the equivalence we presented in Corollary 2.5. In [27], matrices \hat{P} with property ii) of the corollary are called strongly nonsingular. In the case $k = n - k = 1$, we will call such matrices superregular in the next section (see also Remark 3.7). In the same paper [27] it is claimed that the problem of constructing superregular matrices has been solved. Unfortunately this is not true. We will dwell on this in the next section. Therefore, to our knowledge, the problem of constructing strongly nonsingular matrices is still open.

Finally, we wish to verify that for general parameters the properties strongly-MDS, having a maximum distance profile, and m -MDS are not related to each other, meaning that neither of the properties implies the other. We list a suitable example for each case. The column distances have been computed by using straightforward computer routines.

Example 2.12:

- 1) The $(7, 1, 2)$ code over \mathbb{F}_8 given in Example 4.2, item 8) below is strongly-MDS in our sense, but not m -MDS and therefore does not have a maximum distance profile. In this case, $L = 2$ and this is also the memory. It is worth being mentioned that over \mathbb{F}_8 no $(7, 1, 2)$ code with maximum distance profile exists and in particular no systematic m -MDS code with these parameters exists. This can be shown by some lengthy, but straightforward computations.

- 2) The $(4, 1, 2)$ code over \mathbb{F}_{16} with generator matrix

$$\begin{bmatrix} z^2 + \alpha z + 1 & z^2 + \alpha z + \alpha^3 & z^2 + \alpha^6 z + \alpha^3 & \alpha z + \alpha^8 \end{bmatrix}, \\ \text{where } \alpha^4 + \alpha + 1 = 0$$

has a maximum distance profile (thus, by Corollary 2.3, is m -MDS), but is obviously not MDS in our sense and therefore in particular not strongly-MDS.

- 3) The code over \mathbb{F}_{16} with generator matrix

$$\begin{bmatrix} 1 & \alpha^{14} z + \alpha^2 & \alpha^3 z + \alpha^3 \\ \alpha z & \alpha^{11} z + \alpha^8 & \alpha^{10} z + \alpha^{10} \end{bmatrix}$$

is m -MDS but does not have a maximum distance profile.

III. EXISTENCE OF STRONGLY-MDS CODES

During his investigation of algebraic decoding of convolutional codes, Allen conjectured in his dissertation [29] the existence of strongly-MDS convolutional codes in the situation when $k = 1$ and $n = 2$. In this section, we will show the existence of strongly-MDS codes with parameters (n, k, δ) such that $n - k$ divides δ . It follows from (2.11) and Remark 2.10 that these codes also have maximum distance profile. By Theorem 2.6, we have to find an (n, k, δ) code where $(n - k) \mid \delta$ such

that $d_M^c = (n-k)(M+1) + 1$ for $M = \lfloor \frac{\delta}{k} \rfloor + \frac{\delta}{n-k}$. In order to do so, put $m := \frac{\delta}{n-k}$ and let

$$H = [A, B], \text{ where } A = \sum_{i=0}^m A_i D^i \in \mathbb{F}[D]^{(n-k) \times (n-k)}$$

and

$$B = \sum_{i=0}^m B_i D^i \in \mathbb{F}[D]^{(n-k) \times k} \quad (3.1)$$

be a basic parity-check matrix of the desired code. Without loss of generality we may assume $A_0 = I_{n-k}$. The strongly-MDS property can now be expressed as follows.

Theorem 3.1: Let $H \in \mathbb{F}[D]^{n \times (n-k)}$ be as in (3.1), let $A_0 = I_{n-k}$, and define $\mathcal{C} := \{v \in \mathbb{F}((D))^n \mid v\hat{H}^T = 0\}$ to be the code with parity-check matrix H . Furthermore, let

$$A^{-1}B = \sum_{i=0}^{\infty} P_i D^i \in \mathbb{F}((D))^{(n-k) \times k} \quad (3.2)$$

be the Laurent expansion of $A^{-1}B$ over the field $\mathbb{F}((D))$ of Laurent series and for $M := \lfloor \frac{\delta}{k} \rfloor + \frac{\delta}{n-k}$ define

$$\hat{H} := [I_{(M+1)(n-k)} \quad \hat{P}] \quad \text{where} \quad \hat{P} := \begin{bmatrix} P_0 & & & \\ P_1 & P_0 & & \\ \vdots & \vdots & \ddots & \\ P_M & P_{M-1} & \cdots & P_0 \end{bmatrix}. \quad (3.3)$$

We call \hat{H} the M th systematic sliding parity-check matrix of \mathcal{C} . The following conditions are equivalent:

- \mathcal{C} is strongly-MDS, i.e., $d_M^c = (n-k)(M+1) + 1$;
- none of the first k columns of \hat{P} is contained in the span of any other $(M+1)(n-k) - 1$ columns of \hat{H} ;
- each $j \times j$ -submatrix of \hat{P} that does not contain an $s \times t$ -zero block where $s+t \geq j+1$ is nonsingular.

Notice that b) automatically implies that none of the first $n-k$ columns of \hat{H} is in the span of any other $(M+1)(n-k) - 1$ columns.

Proof: After a column permutation the sliding parity-check matrix H_M^c of \mathcal{C} has the form

$$H^c := \left[\begin{array}{ccc|ccc} I & & & B_0 & & \\ A_1 & I & & B_1 & B_0 & \\ \vdots & \vdots & \ddots & \vdots & \vdots & \ddots \\ A_M & A_{M-1} & \cdots & I & B_M & B_{M-1} & \cdots & B_0 \end{array} \right].$$

It is straightforward to see that left multiplication of H^c by the inverse of the first block leads to the matrix \hat{H} of (3.3). After these transformations, Proposition 2.1 applied to the case $j = M$ and $d = (n-k)(M+1) + 1$ translates into the equivalence: \mathcal{C} is strongly-MDS if and only if neither any of the first k columns of H^c nor any of the first $n-k$ columns of the second block of H^c is in the span of any other $(n-k)(M+1) - 1$ columns of \hat{H} . But this, in turn, is equivalent to b) above. The equivalence of a) and c) is obtained from Corollary 2.5. \square

In order to establish the existence of strongly-MDS codes we will proceed as follows. First, we will establish the existence of a systematic sliding parity-check matrix \hat{H} as in (3.3) with property c) of the preceding theorem. Thereafter, we will

show that there exists a basic polynomial matrix $H = [A, B] \in \mathbb{F}[D]^{n \times (n-k)}$ of degree δ such that

$$A^{-1}B = \sum_{i=0}^M P_i D^i + \text{higher powers.}$$

Theorem 3.1 then yields that the code with parity-check matrix H is a strongly-MDS (n, k, δ) code.

As for the first step, let us have a look at the special case of $(2, 1, \delta)$ codes. In this case, $M = 2\delta$ and the systematic sliding parity-check matrix in (3.3) has the form

$$\hat{H} := \left[\begin{array}{c|ccc} 1 & & & h_0 & 0 & \cdots & 0 \\ & 1 & & h_1 & h_0 & \ddots & \vdots \\ & & \ddots & \vdots & \ddots & \ddots & 0 \\ & & & 1 & h_{2\delta} & \cdots & h_1 & h_0 \end{array} \right] \\ =: [I_{2\delta+1}, T] \in \mathbb{F}^{(2\delta+1) \times (4\delta+2)}, \quad \text{where } h_j \in \mathbb{F}. \quad (3.4)$$

As we will see, the existence of matrices T of any given size and the structure above such that \hat{H} has the column property of Theorem 3.1, part b) will be the main tool for the existence of strongly-MDS codes even of length $n > 2$. Therefore, we will concentrate on these matrices first. The main point is to express the column condition on \hat{H} in terms of the minors of T .

Definition 3.2: Let R be a ring. For a matrix $T \in R^{n \times k}$, denote by $T_{j_1, \dots, j_s}^{i_1, \dots, i_r} \in R^{r \times s}$ the $r \times s$ -submatrix obtained from T by picking the rows with indices i_1, \dots, i_r and the columns with indices j_1, \dots, j_s .

In the sequel the following property will play a crucial role.

Definition 3.3: Let \mathbb{F} be field. A lower triangular matrix $T \in \mathbb{F}^{n \times k}$ is said to be superregular,¹ if $T_{j_1, \dots, j_r}^{i_1, \dots, i_r}$ is nonsingular for all $1 \leq r \leq \min\{k, n\}$ and all indices

$$1 \leq i_1 < \cdots < i_r \leq n, \quad 1 \leq j_1 < \cdots < j_r \leq k$$

which satisfy $j_\nu \leq i_\nu$ for $\nu = 1, \dots, r$. We call the submatrices obtained by picking such indices the proper submatrices and their determinants the proper minors of T .

Obviously, a submatrix \hat{T} of T is proper if and only if no diagonal element of \hat{T} comes from strictly above the diagonal of T .

Remark 3.4: Observe that the proper submatrices are the only submatrices which can possibly be nonsingular. This can be seen as follows. If $j_\nu > i_\nu$ for some ν , then in the submatrix $\hat{T} := T_{j_1, \dots, j_r}^{i_1, \dots, i_r}$ the upper right block consisting of the first ν rows and the last $r - \nu + 1$ columns is identically zero. Hence, the first ν rows of \hat{T} can have at most rank $\nu - 1$. In other words, the improper submatrices of T are trivially singular. For example, for $T = (h_{ij})$ we have

$$T_{1,3,4}^{1,2,5} = \begin{bmatrix} h_{11} & 0 & 0 \\ h_{21} & 0 & 0 \\ h_{51} & h_{53} & h_{54} \end{bmatrix}.$$

Before we come to the existence of superregular matrices we will first present the following collection of characterizations of superregular matrices.

¹We adopt this notion from [30], where it has been coined in a slightly different context.

$$T := \begin{bmatrix} (1-\alpha^6)^{-1} & 0 & 0 & \dots & 0 \\ (1-\alpha^5)^{-1} & (1-\alpha^6)^{-1} & 0 & \dots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ (1-\alpha^2)^{-1} & (1-\alpha^3)^{-1} & \dots & (1-\alpha^6)^{-1} & 0 \\ (1-\alpha)^{-1} & (1-\alpha^2)^{-1} & \dots & (1-\alpha^5)^{-1} & (1-\alpha^6)^{-1} \end{bmatrix} \in \mathbb{F}_8^{6 \times 6}$$

Lemma 3.8: Let \mathbb{F} be a field and X_1, \dots, X_l be independent indeterminates over \mathbb{F} . Define the matrix

$$A := \begin{bmatrix} X_1 & 0 & \dots & 0 \\ X_2 & X_1 & \ddots & \vdots \\ \vdots & \ddots & \ddots & 0 \\ X_l & \dots & X_2 & X_1 \end{bmatrix} \in \mathbb{F}(X_1, \dots, X_l)^{l \times l}.$$

Then A is superregular.

Proof: We proceed by contradiction. Assume there exists a singular proper submatrix

$$\hat{A} := A_{j_1, \dots, j_r}^{i_1, \dots, i_r}.$$

We can take the size r to be minimal. Then certainly $r > 1$. By properness, we know that $j_\nu \leq i_\nu$ for $\nu = 1, \dots, r$.

Notice that for $\mu \leq \nu$ the entry of A at the position (ν, μ) is given by $A_\mu^\nu = X_{\nu-\mu+1}$. Hence, the indeterminate with the largest index appearing in \hat{A} is $X_{i_r-j_{r-1}}$. It appears only once in the matrix and that is in the lower left corner. Thus, its coefficient in $\det \hat{A}$ is $\pm \det \tilde{A}$, where

$$\tilde{A} := A_{j_2, \dots, j_r}^{i_1, \dots, i_{r-1}}.$$

Singularity of \hat{A} now implies $\det \tilde{A} = 0$. By minimality of r , this yields that \tilde{A} is an improper submatrix of A , i.e., there exists an index $\tau \in \{2, \dots, r\}$ such that $j_\tau > i_{\tau-1}$. Picking τ minimal we get $i_1 < \dots < i_{\tau-1} < j_\tau < \dots < j_r$ and therefore the first $\tau - 1$ rows of \hat{A} have the form

$$\begin{bmatrix} * & \dots & * & 0 & \dots & 0 \\ \vdots & & \vdots & \vdots & & \vdots \\ * & \dots & * & 0 & \dots & 0 \end{bmatrix}$$

where the block of possibly nonzero elements consists of $\tau - 1$ columns. Hence, \hat{A} is a blocktriangular matrix and we have

$$0 = \det \hat{A} = \det A_{j_1, \dots, j_{\tau-1}}^{i_1, \dots, i_{\tau-1}} \det A_{j_\tau, \dots, j_r}^{i_\tau, \dots, i_r}.$$

Since both factors are proper minors we get a contradiction to the minimality of the size r . \square

The following consequence is standard.

Theorem 3.9: For every $l \in \mathbb{N}$ and every prime number p there exists a finite field \mathbb{F} of characteristic p and a superregular matrix $T \in \mathbb{F}^{l \times l}$ having Toeplitz structure.

Proof: Consider the prime field \mathbb{F}_p and the matrix of the previous lemma with entries in $\mathbb{F}_p(X_1, \dots, X_l)$. All its proper minors are nonzero polynomials in $\mathbb{F}_p[X_1, \dots, X_l]$. Over an algebraic closure $\bar{\mathbb{F}}_p$, a point $a := (a_1, \dots, a_l) \in \bar{\mathbb{F}}_p^l$ can be found such that none of the minors vanishes at a . Hence, the Toeplitz matrix T having $(a_1, \dots, a_l)^T$ as its first column is superregular. Since each a_i is algebraic over \mathbb{F}_p , the matrix T has its entries in a finite field extension \mathbb{F} of \mathbb{F}_p . \square

In particular, for every size $l \in \mathbb{N}$ there exist superregular Toeplitz matrices over a field of characteristic 2. Unfortunately,

the theorem above is nonconstructive and it is not at all clear what the minimum field of characteristic 2 is to allow a superregular Toeplitz matrix of given size $l \times l$. We present some examples.

Example 3.10:

1) Using a computer algebra program one checks that the following matrices are superregular. The first examples are all over prime fields \mathbb{F}_p .

$$\begin{aligned} & \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix} \in \mathbb{F}_2^{2 \times 2}, & \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 2 & 1 & 1 \end{bmatrix} \in \mathbb{F}_3^{3 \times 3} \\ & \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 2 & 1 & 1 & 0 \\ 1 & 2 & 1 & 1 \end{bmatrix} \in \mathbb{F}_5^{4 \times 4}, & \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 2 & 1 & 0 & 0 & 0 \\ 1 & 2 & 1 & 0 & 0 \\ 6 & 1 & 2 & 1 & 0 \\ 4 & 6 & 1 & 2 & 1 \end{bmatrix} \in \mathbb{F}_7^{5 \times 5} \\ & & \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 2 & 1 & 0 & 0 & 0 & 0 \\ 1 & 2 & 1 & 0 & 0 & 0 \\ 1 & 1 & 2 & 1 & 0 & 0 \\ 3 & 1 & 1 & 2 & 1 & 0 \\ 4 & 3 & 1 & 1 & 2 & 1 \end{bmatrix} \in \mathbb{F}_{11}^{6 \times 6} \\ & & \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 7 & 1 & 0 & 0 & 0 & 0 & 0 \\ 13 & 7 & 1 & 0 & 0 & 0 & 0 \\ 2 & 13 & 7 & 1 & 0 & 0 & 0 \\ 1 & 2 & 13 & 7 & 1 & 0 & 0 \\ 4 & 1 & 2 & 13 & 7 & 1 & 0 \\ 14 & 4 & 1 & 2 & 13 & 7 & 1 \end{bmatrix} \in \mathbb{F}_{17}^{7 \times 7}. \end{aligned}$$

The following examples represent superregular matrices over finite fields of characteristic 2. For this assume that α, β , and γ satisfy

$$\alpha^2 + \alpha + 1 = 0, \quad \beta^3 + \beta + 1 = 0, \quad \text{and} \quad \gamma^4 + \gamma + 1 = 0.$$

Then the following matrices are superregular over \mathbb{F}_4 , \mathbb{F}_8 , and \mathbb{F}_{16} , respectively:

$$\begin{aligned} & \begin{bmatrix} 1 \\ \alpha & 1 \\ 1 & \alpha & 1 \end{bmatrix} \in \mathbb{F}_{2^2}^{3 \times 3}, & \begin{bmatrix} 1 & & & & \\ \beta & 1 & & & \\ \beta^3 & \beta & 1 & & \\ \beta & \beta^3 & \beta & 1 & \\ 1 & \beta & \beta^3 & \beta & 1 \end{bmatrix} \in \mathbb{F}_{2^3}^{5 \times 5} \\ & & \begin{bmatrix} 1 & & & & & \\ \gamma & 1 & & & & \\ \gamma^5 & \gamma & 1 & & & \\ \gamma^5 & \gamma^5 & \gamma & 1 & & \\ \gamma & \gamma^5 & \gamma^5 & \gamma & 1 & \\ 1 & \gamma & \gamma^5 & \gamma^5 & \gamma & 1 \end{bmatrix} \in \mathbb{F}_{2^4}^{6 \times 6}. \end{aligned}$$

Assume ϵ, ω satisfy

$$\epsilon^5 + \epsilon^2 + 1 = 0 \quad \text{and} \quad \omega^6 + \omega + 1 = 0.$$

some tedious calculations in the field \mathbb{F}_{64} (where $\omega^6 + \omega + 1 = 0$) one finally arrives at the parity-check matrix

$$H = [1 + \omega^{57}D + \omega^{62}D^2, \omega + \omega^{44}D + \omega^{54}D^2, 1 + \omega^{17}D + \omega^{21}D^2] \in \mathbb{F}_{64}^3.$$

- 3) A strongly-MDS $(4, 3, 1)$ code has free distance 3 and this is identical with the first column distance. It can be obtained from a 6×6 -superregular matrix using the construction of the proof of Theorem 3.11. Indeed, the matrix

$$\hat{H} = \left[\begin{array}{cc|ccc} 1 & 0 & \gamma^5 & \gamma & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & \gamma & \gamma^5 & \gamma^5 & \gamma & 1 \end{array} \right]$$

has been obtained from the superregular Toeplitz matrix of Example 3.10, item 1) and thus it satisfies property b) of Theorem 3.1. Hence, a parity-check matrix of a strongly-MDS $(4, 3, 1)$ code over \mathbb{F}_{16} (where $\gamma^4 + \gamma + 1 = 0$) is given by

$$H = [1, \gamma^5 + D, \gamma + \gamma D, 1 + \gamma^5 D] \in \mathbb{F}_{16}[D]^4.$$

- 4) Of course, not every MDS code is strongly-MDS. For instance, the code with parity-check matrix

$$H = [10 + 3D + 2D^2, 4 + 2D + D^2] \in \mathbb{F}_{11}[D]^2$$

is an MDS code, but not strongly-MDS. In this example, the MDS property follows from the fact that this code is the result of the construction of MDS codes as presented in [12]. However, a $(2, 1, 1)$ code is strongly-MDS if and only if it is an MDS code. This can be checked directly by using Theorem 3.1 and the fact that for the (basic) parity-check matrix $[a_0 + a_1D, b_0 + b_1D]$ of an MDS code all coefficients as well as $a_0b_1 - a_1b_0$ are nonzero.

The next series of examples has been found by completely different methods. They are all cyclic convolutional codes in the sense of [13], [35], [14], [15]. In those papers, convolutional codes having some additional algebraic structure are being investigated. This additional structure is a generalization of cyclicity of block codes but is a far more complex notion for convolutional codes. In particular, cyclicity of convolutional codes does *not* mean invariance under the cyclic shift in $\mathbb{F}[D]^n$. We will not go into the details but rather refer to [13], [35]. However, in order to understand and test the following examples there is no need in understanding the concept of cyclicity for convolutional codes since in the following we provide all information needed to specify the codes. We present the generator matrices and also provide all column distances; they have been computed with a computer algebra program. All matrices given below are minimal basic. We would like to mention that just like for cyclic block codes, the length of the code and the characteristic of the field have to be coprime. Therefore, only codes with odd length are given below.

One should note that most of the following codes exist over comparatively smaller alphabets than the examples of Example 4.1. However, we do not know any general construction

for strongly-MDS cyclic convolutional codes yet. But the abundance of (small) examples suggests that such a construction might be possible and might lead to smaller alphabets for given parameters than the construction of the last section. We will leave this as an open question for future research.

Example 4.2:

- 1) A strongly-MDS $(3, 1, 1)$ code over \mathbb{F}_4

$$G = [\alpha + \alpha D, \alpha^2 + \alpha D, 1 + \alpha D].$$

The column distances are $d_0^c = 3$, $d_1^c = 5$, $d_j^c = 6$ for $j \geq 2$.

- 2) A strongly-MDS $(3, 1, 2)$ code over \mathbb{F}_{16} (where $\beta^4 + \beta + 1 = 0$)

$$G = [\beta + \beta D + D^2, \beta^6 + \beta D + \beta^{10}D^2, \beta^{11} + \beta D + \beta^5D^2].$$

The column distances are $d_0^c = 3$, $d_1^c = 5$, $d_2^c = 7$, $d_j^c = 9$ for $j \geq 3$.

- 3) A strongly-MDS $(3, 2, 2)$ code over \mathbb{F}_{16}

$$G = \begin{bmatrix} \beta^5 + \beta^4D & \beta^3 + \beta^8D & \beta^9 + \beta^2D \\ \beta^9 + \beta^{12}D & \beta^5 + \beta^{14}D & \beta^3 + \beta^3D \end{bmatrix}.$$

The column distances are $d_0^c = 2$, $d_1^c = 3$, $d_2^c = 4$, $d_j^c = 5$ for $j \geq 3$.

- 4) A strongly-MDS $(5, 1, 1)$ code over \mathbb{F}_{16}

$$G = [\beta + \beta D, \beta^{13} + \beta^{10}D, \beta^{10} + \beta^4D, \beta^7 + \beta^{13}D, \beta^4 + \beta^7D].$$

The column distances are $d_0^c = 5$, $d_1^c = 9$, $d_j^c = 10$ for $j \geq 2$.

- 5) A strongly-MDS $(5, 1, 2)$ code over \mathbb{F}_{16}

$$G = [\beta + \beta^4D + \beta D^2, \beta^7 + \beta D + \beta^{10}D^2, \beta^{13} + \beta^{13}D + \beta^4D^2, \beta^4 + \beta^{10}D + \beta^{13}D^2, \beta^{10} + \beta^7D + \beta^7D^2].$$

The column distances are $d_0^c = 5$, $d_1^c = 9$, $d_2^c = 13$, $d_j^c = 15$ for $j \geq 3$.

- 6) A strongly-MDS $(5, 2, 2)$ code over \mathbb{F}_{16} as shown at the bottom of the page. The column distances are $d_0^c = 4$, $d_1^c = 7$, $d_j^c = 9$ for $j \geq 2$.

- 7) A strongly-MDS $(7, 1, 1)$ code over \mathbb{F}_8 (where $\gamma^3 + \gamma + 1 = 0$)

$$G = [\gamma + \gamma D, \gamma^3 + D, \gamma^5 + \gamma^6D, 1 + \gamma^5D, \gamma^2 + \gamma^4D, \gamma^4 + \gamma^3D, \gamma^6 + \gamma^2D].$$

The column distances are $d_0^c = 7$, $d_1^c = 13$, $d_j^c = 14$ for $j \geq 2$.

- 8) A strongly-MDS $(7, 1, 2)$ code over \mathbb{F}_8

$$G = [\gamma^2 + \gamma D + D^2, \gamma^5 + \gamma^3D + \gamma^6D^2, \gamma + \gamma^5D + \gamma^5D^2, \gamma^4 + D + \gamma^4D^2, 1 + \gamma^2D + \gamma^3D^2, \gamma^3 + \gamma^4D + \gamma^2D^2, \gamma^6 + \gamma^6D + \gamma D^2].$$

The column distances are $d_0^c = 7$, $d_1^c = 13$, $d_2^c = 18$, $d_j^c = 21$ for $j \geq 3$.

$$G = \begin{bmatrix} \beta + \beta D & \beta^{13} + \beta^{10}D & \beta^{10} + \beta^4D & \beta^7 + \beta^{13}D & \beta^4 + \beta^7D \\ 1 + \beta^5D & \beta^3 + \beta^{11}D & \beta^6 + \beta^2D & \beta^9 + \beta^8D & \beta^{12} + \beta^{14}D \end{bmatrix}.$$

- 9) It is worth noting that there does not exist even an MDS $(7, 2, 2)$ code over \mathbb{F}_8 , since the generalized Singleton bound in this case is 13, but due to the Griesmer bound (see [1, p. 133] for the binary case) the parameters of an (n, k, δ) code over \mathbb{F}_q with memory m and distance d satisfy

$$\sum_{l=0}^{k(m+i)-\delta-1} \left\lceil \frac{d}{q^l} \right\rceil \leq n(m+i), \quad \text{for all } i \in \mathbb{N}_0.$$

Hence, a $(7, 2, 2)$ code over \mathbb{F}_8 with memory 1 has at most distance 12. The inequality applied to $i = 1$ shows that the field size has to be at least 13 in order to allow the existence of an MDS $(7, 2, 2)$ code.

One should notice that the codes in Example 4.2, items 1)–7) are not only strongly-MDS but also have *all* column distances being optimal in the sense that they reach the upper bound given in Proposition 2.2. In particular, they also have a maximum distance profile in the sense of Definition 2.9. For the $(7, 1, 2)$ code in item 8), only the second column distance is not optimal, but rather one less than the upper bound, which is 19 in this case. The implications of this have been discussed already in Example 2.12, item 1).

V. THE DUAL OF A STRONGLY MDS-CODE

In this section, we will present some results concerning the dual code of a strongly-MDS code. The main result shows that a convolutional code has a maximum distance profile if and only if its dual has this property. This then implies for certain parameters that a code is strongly-MDS if and only if its dual has this property. These results are very appealing as they generalize the situation for block codes.

Recall that if

$$\mathcal{C} = \{uG \mid u \in \mathbb{F}[D]^k\} = \{v \in \mathbb{F}[D]^n \mid vH^\top = 0\} \subseteq \mathbb{F}[D]^n$$

is an (n, k, δ) code with generator matrix $G \in \mathbb{F}[D]^{k \times n}$ and parity-check matrix $H \in \mathbb{F}[D]^{(n-k) \times n}$, then the dual code, defined as

$$\mathcal{C}^\perp = \{w \in \mathbb{F}[D]^n \mid wv^\top = 0 \text{ for all } v \in \mathcal{C}\}$$

is given by

$$\mathcal{C}^\perp = \{uH \mid u \in \mathbb{F}[D]^{n-k}\} = \{w \in \mathbb{F}[D]^n \mid wG^\top = 0\}$$

and thus an $(n, n-k, \delta)$ code. In contrast to the block code situation, almost nothing is known about the relation between the distances of a code and its dual. In particular, it has been shown in [36] that no MacWilliams identity relating the weight distributions of \mathcal{C} and \mathcal{C}^\perp exists. In block code theory, a very simple relation between the distances of a code and its dual is given in the case of MDS codes. In fact, if \mathcal{C} is an MDS (n, k) block code, then the dual \mathcal{C}^\perp is an MDS $(n, n-k)$ code, see [33, Ch. 11, Par. 2] and very specific knowledge on the weight enumerator and its dual is known [33, Ch. 11]. Therefore, it is quite natural to investigate whether the dual of an MDS (or strongly-MDS) convolutional code is MDS (or strongly-MDS), too. Unfortunately, this is in general not the case.

Example 5.1: In general, the dual of a strongly-MDS code is not even an MDS code. This can be seen from the dual of the code given in Example 4.1, item 3). The dual has generator matrix

$$G = [1, \gamma^5 + D, \gamma + \gamma D, 1 + \gamma^5 D] \in \mathbb{F}_{16}[D]^4$$

which obviously has weight less than the generalized Singleton bound 8 (see Theorem 2.6).

As we will show next, the property of maximum distance profile carries over under dualization. In addition, for specific code parameters the strongly-MDS property carries over to the dual code as well. To this end, recall from Definition 2.8 that an (n, k, δ) code is strongly-MDS if the M th column distance attains the generalized Singleton bound where

$$M = \left\lfloor \frac{\delta}{k} \right\rfloor + \left\lceil \frac{\delta}{n-k} \right\rceil.$$

Thus, the dual code \mathcal{C}^\perp is MDS if the \hat{M} th column distance attains the generalized Singleton bound where

$$\hat{M} = \left\lfloor \frac{\delta}{n-k} \right\rfloor + \left\lceil \frac{\delta}{k} \right\rceil.$$

Obviously, these two numbers differ by one when k divides δ but $n-k$ does not or *vice versa*. What remains equal for both the code and its dual is the quantity

$$L = \left\lfloor \frac{\delta}{k} \right\rfloor + \left\lfloor \frac{\delta}{n-k} \right\rfloor$$

used in Definition 2.9, where we introduced the concept of maximum distance profile.

Theorem 2.4 provides us with the following nice duality result.

Theorem 5.2: An (n, k, δ) code $\mathcal{C} \subseteq \mathbb{F}[D]^n$ has a maximum distance profile if and only if the dual code $\mathcal{C}^\perp \subseteq \mathbb{F}[D]^n$ has this property.

Proof: Let \mathcal{C} have generator matrix G and parity-check matrix H as given in (2.1) and (2.2). Assume \mathcal{C} has a maximum distance profile. By Theorem 2.4, every $(L+1)(n-k) \times (L+1)(n-k)$ full-size minor formed from the columns of H_L^\perp with indices $1 \leq r_1 < \dots < r_{(L+1)(n-k)}$, where $r_{s(n-k)} \leq sn$ for $s = 1, \dots, L$, is nonzero.

Consider now the dual code \mathcal{C}^\perp which is defined as the row-space of the $(n-k) \times n$ matrix H . It follows from (2.4) that the L th column distance of the dual code \mathcal{C}^\perp is given by

$$\hat{d}_L^\perp = \min \{ \text{wt}((u_L, \dots, u_0)H_L^\perp) \mid u_i \in \mathbb{F}^{n-k}, u_0 \neq 0 \}.$$

Taking the reversed ordering into account, we obtain again from Theorem 2.4 that the dual code \mathcal{C}^\perp has maximum distance profile as well. \square

Corollary 5.3: When both k and $n-k$ divide δ then an (n, k, δ) code $\mathcal{C} \subseteq \mathbb{F}[D]^n$ is strongly-MDS if and only if $\mathcal{C}^\perp \subseteq \mathbb{F}[D]^n$ has this property.

Proof: From $k \mid \delta$ and $(n-k) \mid \delta$, it follows that $L = M$ and

$$d_M^\perp = (n-k) \left(\frac{\delta}{k} + 1 \right) + \delta + 1$$

which is the generalized Singleton bound of the code \mathcal{C} and

$$\hat{d}_M^{\mathcal{C}} = k \left(\frac{\delta}{n-k} + 1 \right) + \delta + 1$$

which is the generalized Singleton bound of the dual code \mathcal{C}^\perp . \square

The preceding result gives us another class of strongly-MDS codes by dualizing Theorem 3.11.

Corollary 5.4: For every $n, \delta \in \mathbb{N}_0$ such that both k and $n - k$ divide δ and every prime number p there exists a strongly-MDS (n, k, δ) code over some suitably large field of characteristic p .

Example 5.5:

- a) Corollary 5.3 tells us that the duals of the $(2, 1, \delta)$ codes given in Example 4.1, item 1) are strongly-MDS. But this is obviously so, since they are—up to ordering—identical to the given codes.

- b) Dualizing the code of Example 4.1, item 2) gives us a strongly-MDS $(3, 1, 2)$ code with generator matrix

$$G = [1 + \omega^{57}D + \omega^{62}D^2, \omega + \omega^{44}D + \omega^{54}D^2, 1 + \omega^{17}D + \omega^{21}D^2] \in \mathbb{F}_{64}^3.$$

- c) Dualizing the codes given in Example 4.2, items 2) and 3), we obtain another two strongly-MDS codes with generator matrices

$$H_1 = \begin{bmatrix} 1 & \beta D + \beta^9 & \beta^6 D + \beta^8 \\ \beta^{14} D & \beta^7 D + \beta^6 & \beta^8 D + \beta \end{bmatrix} \in \mathbb{F}_{16}^{2 \times 3}$$

and

$$H_2 = [D^2 + D + \beta^2, \beta^{10} D^2 + D + \beta^7, \beta^5 D^2 + D + \beta^{12}] \in \mathbb{F}_{16}^3.$$

It is known that these codes are also cyclic convolutional codes in the sense of [13], see [13, Theorem 7.5].

Finally, we would like to mention that even in the case where $k \mid \delta$ and $(n - k) \mid \delta$, the dual of an MDS code is not MDS in general. An example is given by the following code.

Example 5.6: The $(3, 1, 2)$ code $\mathcal{C} \subseteq \mathbb{F}[D]^3$, where $\mathbb{F} = \mathbb{F}_{16}$, with generator matrix

$$G = [1 + \beta D + \beta^4 D^2, \beta^{10} + \beta^2 D + \beta^4 D^2, \beta^8 + \beta^5 D + D^2]$$

and parity-check matrix

$$H = \begin{bmatrix} 1 & \beta^{14} D + \beta^2 & \beta^3 D + \beta^3 \\ \beta D & \beta^{11} D + \beta^8 & \beta^{10} D + \beta^{10} \end{bmatrix}$$

is an MDS code, but not strongly-MDS. It satisfies $d_3^{\mathcal{C}} = 8$ and $d_4^{\mathcal{C}} = 9$. The dual code generated by H is not MDS. Its distance is 4.

VI. ESTIMATES FOR THE EXTENDED ROW DISTANCES

In this section, we will use the information about the column distances in order to present a lower bound for the extended row distances of a strongly-MDS code with unit memory.

For this, let $G = G_0 + G_1 z \in \mathbb{F}[z]^{k \times n}$ be the generator matrix of a unit memory code of degree k , thus, G_1 has full row rank, and let the code be strongly-MDS and have a maximum

distance profile. Hence, $d_j^{\mathcal{C}} = (n - k)(j + 1) + 1$ for all $j = 0, \dots, M - 1$ where $M := 1 + \lceil \frac{k}{n-k} \rceil$ and we have

$$d_M^{\mathcal{C}} = 2n - k + 1 = d_{\text{free}}(\text{im } G).$$

Denote by \hat{d}_j^r the j th extended row distances of the code, thus, \hat{d}_j^r is the minimum weight of all codewords $v = \sum_{i=0}^{j-1} u_i z^i G$ of degree j where $u_i \neq 0$ for all $i = 0, \dots, j - 1$. Define

$$A_M^{\mathcal{C}} = \begin{bmatrix} G_1 & & & \\ G_0 & G_1 & & \\ & G_0 & \ddots & \\ & & \ddots & G_1 \\ & & & & G_0 \end{bmatrix} \in \mathbb{F}^{(M+1)k \times Mn}.$$

Lemma 6.1: Let $u_0 \neq 0$. Then

$$\text{wt}((u_0, u_1, \dots, u_M) A_M^{\mathcal{C}}) \geq n - k + 1.$$

Proof: Let $(u_0, u_1, \dots, u_M) A_M^{\mathcal{C}} = (v_1, \dots, v_M)$. Then $(u_0, u_1, \dots, u_M) G_M^{\mathcal{C}} = (v_0, v_1, \dots, v_M)$ for some $v_0 \in \mathbb{F}^n$. Since $u_0 \neq 0$, we have

$$\text{wt}(v_0, v_1, \dots, v_M) \geq d_M^{\mathcal{C}} = 2n - k + 1.$$

Estimating the weight of v_0 by n we obtain the desired result. \square

Theorem 6.2: Let $j \geq M$ and write $j = aM + t$ where $a \in \mathbb{N}$ and $0 \leq t < M$. Then we have

$$\hat{d}_j^r \geq \frac{n - k + 1}{M} j + (n - k)(M - 1) + \max\{0, t(n - k) + 1 - n\}.$$

Hence, the extended row distances are bounded from below by a linear function with slope $\frac{n-k+1}{M}$.

Note that the constant part of this linear function is always positive.

Proof: Write the matrix at the top of the following page. We have to estimate $\text{wt}((u_0, \dots, u_j) G_j^r)$ where all $u_i \neq 0$. Thus, we may use the lower bound $d_{M-1}^{\mathcal{C}} = M(n - k) + 1$ for the first block, the lemma for the next $a - 1$ blocks, $\max\{0, d_{t-1}^{\mathcal{C}} - n\}$ and $n - k + 1$ for the last two blocks, respectively. Hence,

$$\begin{aligned} \hat{d}_j^r &\geq M(n - k) + 1 + (a - 1)(n - k + 1) \\ &\quad + \max\{0, t(n - k) + 1 - n\} + n - k + 1 \\ &= (n - k)(M + a) + a + 1 + \max\{0, t(n - k) + 1 - n\} \\ &= (n - k) \left(M + \left\lfloor \frac{j}{M} \right\rfloor \right) + \left\lfloor \frac{j}{M} \right\rfloor + 1 + \max\{0, t(n - k) + 1 - n\} \\ &\geq (n - k) \left(\frac{j}{M} + M - 1 \right) + \frac{j}{M} + \max\{0, t(n - k) + 1 - n\} \\ &= \frac{n - k + 1}{M} j + (n - k)(M - 1) + \max\{0, t(n - k) + 1 - n\}. \quad \square \end{aligned}$$

Remark 6.3: We computed the weight distribution for some of our codes (those with number of states not bigger than 64), and in all cases we even obtain $\hat{d}_j^r = (n - k)(j + 1) + 2$. This is in general a much better slope than the estimate of the preceding theorem.

VII. CONCLUSION

In this paper, we introduced two new classes of convolutional codes called strongly-MDS convolutional codes and codes

$$G_j^r = \left(\begin{array}{c|c|c|c|c}
 \begin{array}{ccc} G_0 & G_1 & \\ & G_0 & \ddots \\ & & \ddots & G_1 \\ & & & G_0 \end{array} & & & & \\
 \hline
 & \begin{array}{ccc} G_1 & & \\ G_0 & G_1 & \\ & G_0 & \ddots \\ & & \ddots & G_1 \\ & & & G_0 \end{array} & & & \\
 \hline
 & & & & \\
 \hline
 & & & \ddots & \\
 \hline
 & & & & G_1 \\
 \hline
 & & & \begin{array}{ccc} G_0 & G_1 & \\ & G_0 & \ddots \\ & & \ddots & G_1 \\ & & & G_0 \end{array} & G_1 \\
 \hline
 \end{array} \right) \in \mathbb{F}^{jk \times (j+1)n}.$$

$\underbrace{\hspace{10em}}_{M \text{ blocks}} \quad \underbrace{\hspace{10em}}_{M \text{ blocks}} \quad \dots \quad \underbrace{\hspace{10em}}_{t \text{ blocks}} \quad \underbrace{\hspace{10em}}_{1 \text{ block}}$
 $a \text{ times}$

having maximum distance profile. Strongly-MDS convolutional codes have the property that the generalized Singleton bound is attained at the earliest possible column distance. Codes with maximum distance profile have a maximal possible increase of the column distances.

From an applications point of view, strongly-MDS convolutional codes are particularly suited in situations where codes over large alphabets are required and in situations where algebraic decoding is desirable. Hadjicostis [21], [22] has recently demonstrated that convolutional codes over large alphabets are very desirable in areas of process control via linear finite-state machines, where large numbers of nonconcurrent errors should be detected and corrected. It seems that strongly-MDS convolutional codes have potential for such applications.

APPENDIX A

We will need the following lemma.

Lemma A.1: Let $A \in \mathbb{F}^{k \times n}$ and $B \in \mathbb{F}^{n \times (n-k)}$ such that

$$AB = 0 \text{ and } \text{rank } A = k, \quad \text{rank } B = n - k.$$

Then the following are equivalent:

- a) the $k \times k$ -submatrix of A consisting of the columns with indices $1 \leq t_1 < \dots < t_k \leq n$ is singular;
- b) the $(n-k) \times (n-k)$ -submatrix of B obtained by taking the rows with indices in $\{1, \dots, n\} \setminus \{t_1, \dots, t_k\}$ is singular;

Proof: Without loss of generality, assume $(t_1, \dots, t_k) = (1, \dots, k)$ and partition $A = (A_1 \ A_2)$, where A_1 is the $k \times k$ submatrix under consideration. If A_1 is invertible then

$$\ker A = \text{colspan}_{\mathbb{F}} \begin{pmatrix} A_1^{-1} A_2 \\ -I_{n-k} \end{pmatrix} = \text{colspan}_{\mathbb{F}}(B).$$

This shows that the bottom $(n-k) \times (n-k)$ -submatrix of B is invertible. \square

Proof of Theorem 2.4:

i) \Rightarrow ii): Assume there are indices $1 \leq t_1 < \dots < t_{(j+1)k}$ satisfying $t_{sk+1} > sn$ for $s = 1, \dots, j$ whose corresponding minor is zero. It follows that there is a vector $u = (u_0, \dots, u_j)$ such that uG_j^c has zero coordinates at positions $t_1, \dots, t_{(j+1)k}$. Let $\ell := \min\{i \mid u_i \neq 0\}$. Consider the vector

$$(u_\ell, \dots, u_j)G_{j-\ell}^c \in \mathbb{F}^{(j-\ell+1)n}.$$

The weight of this vector is at most $(j-\ell+1)(n-k)$ as there are at least $(j-\ell+1)k$ zero coordinates. From (2.4), it follows that $d_{j-\ell}^c \leq (j-\ell+1)(n-k)$ and by Corollary 2.3 we also have $d_j^c \leq (j+1)(n-k)$, contradicting i).

ii) \Rightarrow i): Assume that $d_j^c \leq (n-k)(j+1)$. Let

$$m := \min\{i \mid d_i^c \leq (n-k)(i+1)\}.$$

It follows that there is a vector $u = (u_0, \dots, u_m)$, $u_0 \neq 0$ such that uG_m^c has at least $k(m+1)$ zeros. As a submatrix inside G_j^c we select the columns corresponding to the first $k(m+1)$ positions where uG_m^c has a zero and we augment it by the last $k(j-m)$ columns of G_j^c . We call the indices of the selected columns $t_1, \dots, t_{(j+1)k}$. This gives a $(j+1)k \times (j+1)k$ full-size minor and we claim that this minor is zero and that the indices $t_1, \dots, t_{(j+1)k}$ satisfy $t_{sk+1} > sn$ for $s = 1, \dots, j$. In order to prove the latter, note that $d_i^c = (n-k)(i+1) + 1$ for $i = 0, \dots, m-1$. It therefore follows that $(u_0, \dots, u_i)G_i^c$ has at most $k(i+1) - 1$ zeros for $i = 0, \dots, m-1$. In particular, $t_{sk+1} > sn$ for $s = 1, \dots, m$. Clearly, it is also true for $s = m+1, \dots, j$. It remains to be shown that the minor is zero. For this, note that the selected matrix has the form $\begin{bmatrix} A & B \\ 0 & C \end{bmatrix}$, where A is an $(m+1)k \times (m+1)k$ submatrix of G_m^c which is singular by construction. The full size minor is therefore zero.

on the first two rows, thus, $i_2 = i_1 + 1$ and $j_1 = i_1$. Since $j_2 \in \{i_2, i_2 - 1\} = \{i_1 + 1, i_1\}$ and $j_2 > j_1$, we can only have $j_2 = i_1 + 1$. Then the first row will have only one nonzero entry equal to 1 on the first position, and applying cofactor expansion along that row, we obtain again (B3) and thus $\det \hat{X} > 0$. We now proceed by induction on k in order to prove the desired result for X^k where $k > 1$. Assume X^{k-1} has the stated property. Using $X^k = X \cdot X^{k-1}$ and the Cauchy–Binet formula for minors we obtain

$$\det \hat{X} = \sum_{\substack{1 \leq s_1 < \dots < s_r \leq n, \\ s_l \in \{i_l, i_l - 1\} \cap \{j_l, j_l + 1, \dots, j_l + k - 1\}}} \det X_{s_1, \dots, s_r}^{i_1, \dots, i_r} \cdot \det (X^{k-1})_{j_1, \dots, j_r}^{s_1, \dots, s_r}.$$

Due to part 1) of the proof, the sum indeed expands only over the given indices. By induction, all nonsingular submatrices of both matrices X and X^{k-1} have positive determinant, hence if there are any nonzero terms in the sum, it is necessarily positive. Therefore, the only thing left to be proven is that there is a nonzero term in the above sum. But all products of the form $\det X_{i_1, \dots, i_r}^{i_1, \dots, i_r} \cdot \det (X^{k-1})_{j_1, \dots, j_r}^{i_1, \dots, i_r}$ with

$$j_l \in \{i_l, i_l - 1, i_l - 2, \dots, i_l - (k - 1)\}$$

for all l are nonzero. Thus, $\det \hat{X} > 0$ and the proof is complete. \square

APPENDIX C

Proof of Theorem 3.11:

Step 1: We will show the existence of a matrix \hat{P} as in (3.3) satisfying part c) of Theorem 3.1. This can be accomplished as follows. Define $\tau = (M+1)(n-1)$ and pick a superregular Toeplitz matrix

$$T = \begin{bmatrix} t_0 & 0 & \dots & 0 \\ t_1 & t_0 & \ddots & \vdots \\ \vdots & \ddots & \ddots & 0 \\ t_{\tau-1} & \dots & t_1 & t_0 \end{bmatrix} \in \mathbb{F}^{\tau \times \tau}.$$

Theorem 3.9 guarantees the existence of such a matrix over a suitably large field of characteristic p . For $l = 0, \dots, M$ define the matrices

$$P_l = T_{1, \dots, k}^{l(n-1)+k, l(n-1)+k+1, \dots, (l+1)(n-1)} \in \mathbb{F}^{(n-k) \times k}.$$

Then, due to the Toeplitz structure of T , we have for all $r = 0, \dots, M$

$$P_l = T_{r(n-1)+1, \dots, r(n-1)+k}^{(l+r)(n-1)+k, \dots, (l+r+1)(n-1)} \quad \text{for } l = 0, \dots, M - r$$

and therefore the matrix \hat{P} in (3.3) is obtained from T by picking the rows with indices

$$k, k+1, \dots, n-1, n-1+k, \dots, 2(n-1), \dots, M(n-1)+k, \dots, (M+1)(n-1)$$

and the columns with indices

$$1, \dots, k, n, \dots, n-1+k, 2(n-1)+1, \dots, 2(n-1)+k, \dots, M(n-1)+1, \dots, M(n-1)+k.$$

But then it is obvious that the matrix \hat{P} inherits from T the property that all $j \times j$ -submatrices not containing an $s \times r$ -zero block where $s+r \geq j+1$ are nonsingular, cf. Remark 3.7. This provides us with a matrix \hat{H} as in (3.3) satisfying the equivalent conditions of Theorem 3.1.

Step 2: We now establish the existence of a polynomial matrix $H \in \mathbb{F}[D]^{(n-k) \times n}$ having \hat{H} in (3.3) as M th systematic sliding parity-check matrix. For this, define $m := \frac{\delta}{n-k}$ and start with H as in (3.1). Without loss of generality we may assume $A_0 = I_{n-k}$. Then (3.2) tells us that we need matrices A and B satisfying

$$B = A \left(\sum_{i=0}^M P_i D^i + \text{higher terms} \right). \quad (\text{C1})$$

Comparing the coefficients of D^{m+1}, \dots, D^M we obtain the matrix equation

$$[A_m \quad \dots \quad A_1] \begin{bmatrix} P_{M-m} & \dots & P_1 \\ P_{M-m+1} & \dots & P_2 \\ \vdots & & \vdots \\ P_{M-1} & \dots & P_m \end{bmatrix} = -[P_M \quad \dots \quad P_{m+1}]. \quad (\text{C2})$$

Denote the matrix occurring on the left-hand side by \mathcal{P} . In order to see that the matrix equation is solvable we will show that \mathcal{P} has full column rank. Notice that $\mathcal{P} \in \mathbb{F}^{m(n-k) \times (M-m)k}$. Using $M = \lfloor \frac{\delta}{k} \rfloor + \frac{\delta}{n-k}$ and $m = \frac{\delta}{n-k}$ one can easily see that $(M-m)k \leq m(n-k)$. But then the full column rank of \mathcal{P} follows from part c) of Theorem 3.1. Thus, we can find matrices A_1, \dots, A_m satisfying (C2). Comparing now the powers of D^0, \dots, D^m in (C1) we obtain B_0, \dots, B_m . Then the equation is fully satisfied by setting P_{M+1}, P_{M+2}, \dots suitably.

Step 3: It remains to see that the code

$$\mathcal{C} := \{v \in \mathbb{F}[D]^n \mid vH^T = 0\}$$

has degree δ . But this follows directly from the construction. Indeed, recall that $M = \lfloor \frac{\delta}{k} \rfloor + \frac{\delta}{n-k}$. Therefore, using Theorem 3.1, part a), we know that the M th column distance of \mathcal{C} satisfies

$$\begin{aligned} d_M^c &= (n-k)(M+1) + 1 \\ &= (n-k) \left(\lfloor \frac{\delta}{k} \rfloor + 1 \right) + \delta + 1. \end{aligned}$$

Now the generalized Singleton bound in Theorem 2.6 shows that the code \mathcal{C} cannot have a degree smaller than δ . Thus, \mathcal{C} is a strongly-MDS (n, k, δ) code, $n-k \mid \delta$ and the proof is complete. \square

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